TOPIC

LATTICE-BASED ACCESS-CONTROL MODELS

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LATTICE-BASED MODELS

- Denning's axioms
- Bell-LaPadula model (BLP)
- Biba model and its duality (or equivalence) to BLP
- Dynamic labels in BLP

DENNING'S AXIOMS

$$<$$
 SC, \rightarrow , \oplus $>$

SC

 \rightarrow \subseteq SC X SC

⊕: SC X SC -> SC

set of security classes

flow relation (i.e., can-flow)

class-combining operator

DENNING'S AXIOMS

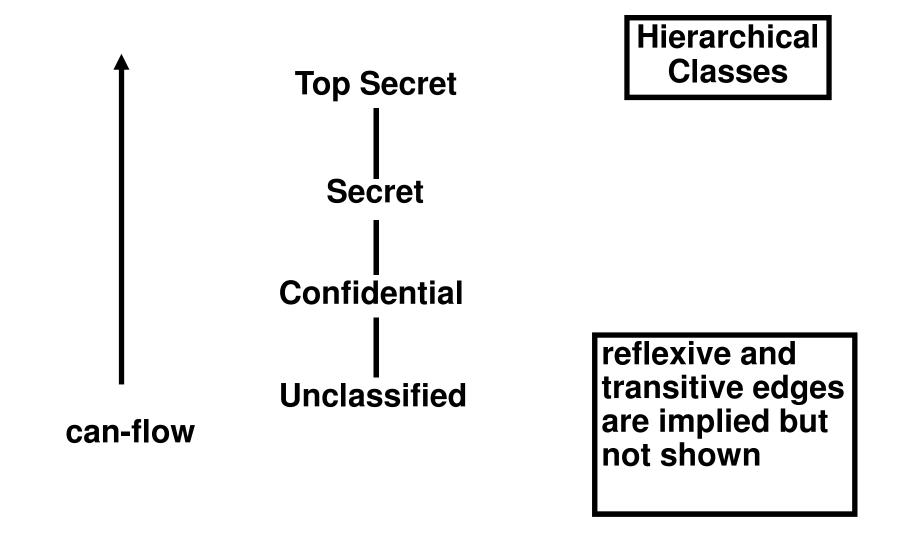
$$<$$
 SC, \rightarrow , \oplus $>$

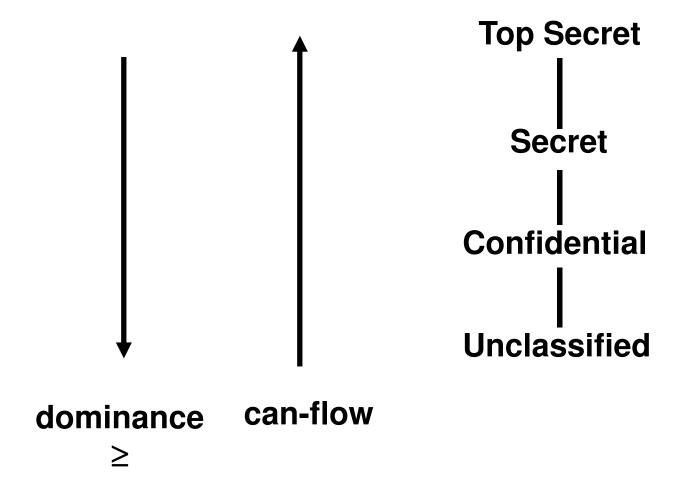
- 1 SC is finite
- $2 \rightarrow is a partial order on SC$
- 3 SC has a lower bound L such that L \rightarrow A for all A \in SC
- 4 ⊕ is a least upper bound (lub) operator on SC

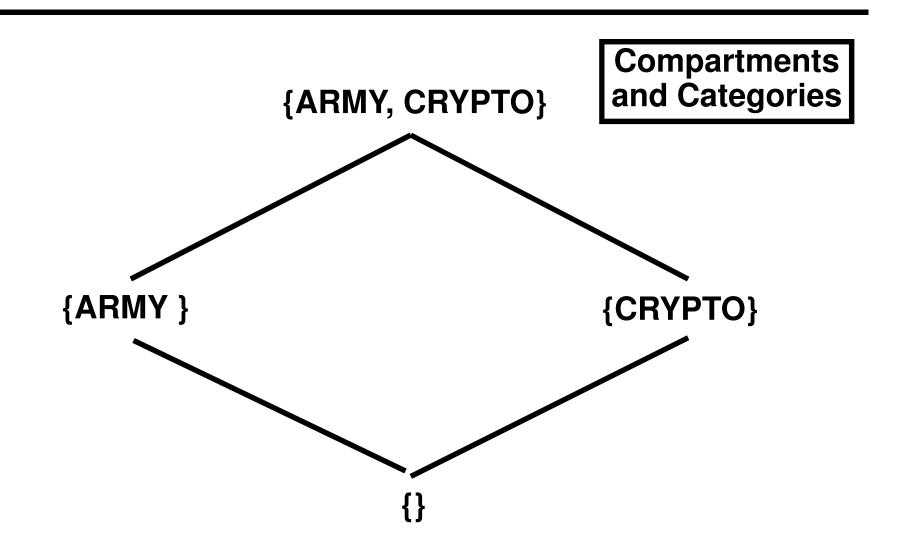
Justification for 1 and 2 is stronger than for 3 and 4. In practice we may therefore end up with a partially ordered set (poset) rather than a lattice.

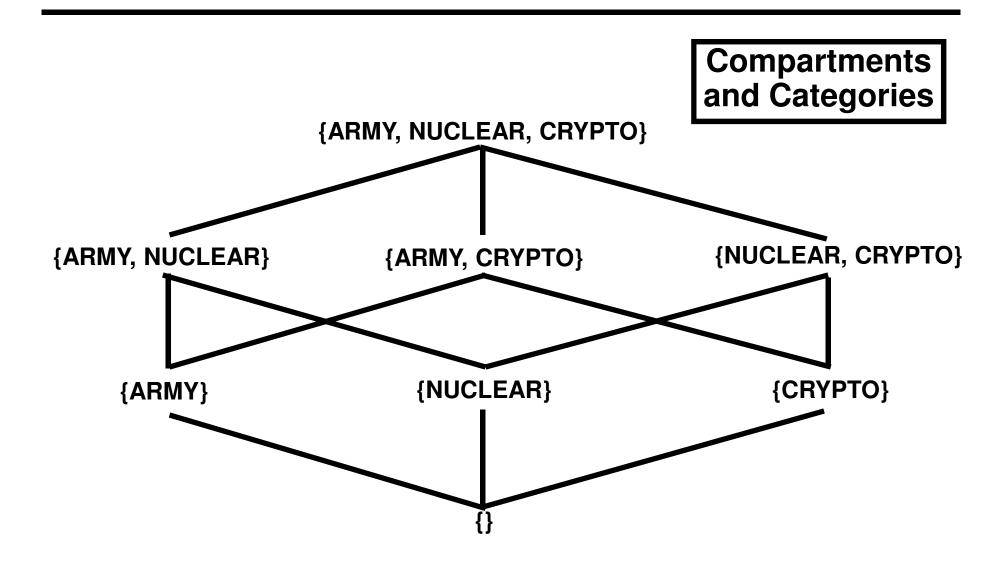
DENNING'S AXIOMS IMPLY

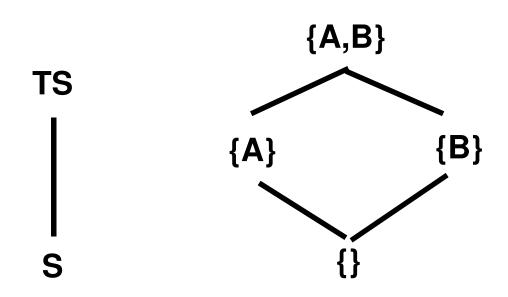
- SC is a universally bounded lattice
- there exists a Greatest Lower Bound (glb) operator ⊗ (also called meet)
- there exists a highest security class H





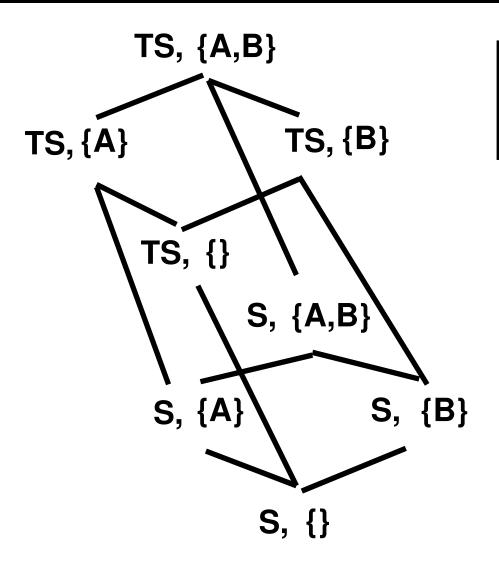






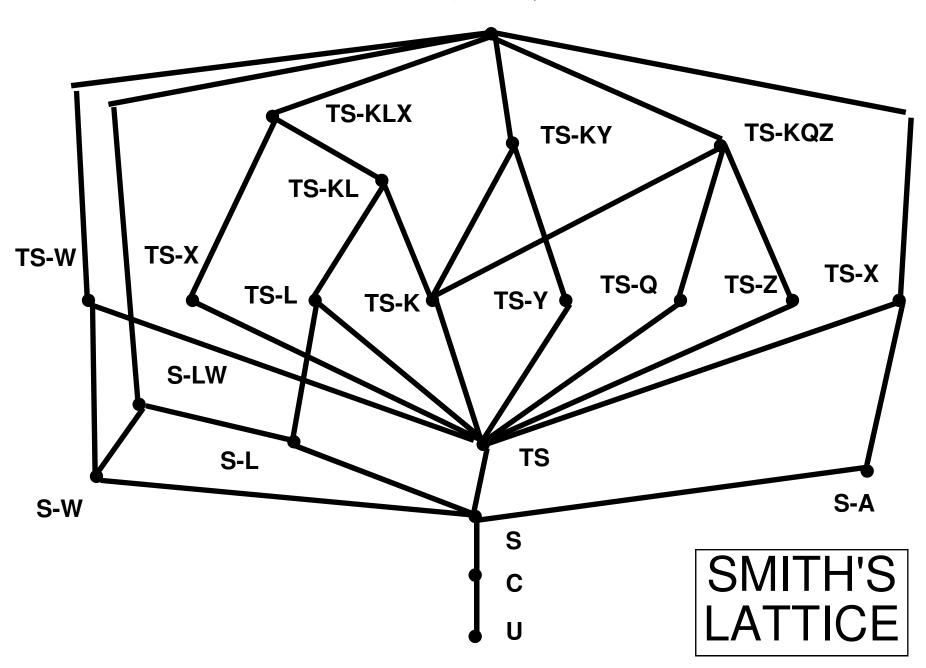
Hierarchical Classes with Compartments

product of 2 lattices is a lattice



Hierarchical Classes with Compartments

TS-AKLQWXYZ



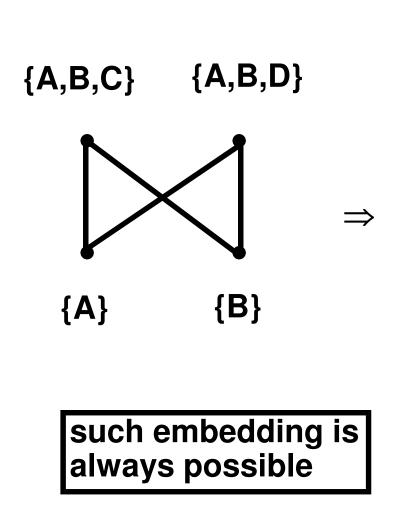
SMITH'S LATTICE

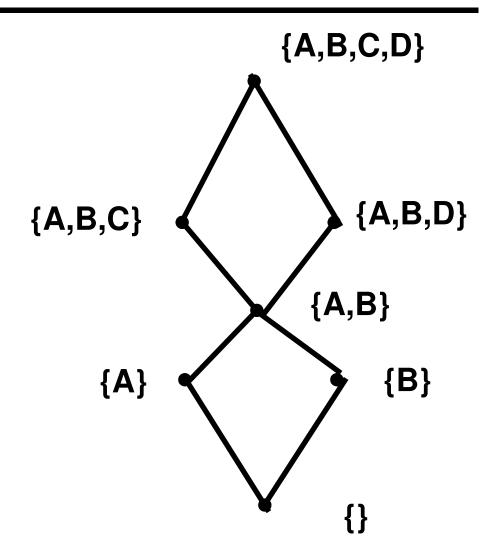
- With large lattices a vanishingly small fraction of the labels will actually be used
 - Smith's lattice: 4 hierarchical levels, 8 compartments, therefore number of possible labels = 4*2^8 = 1024
 Only 21 labels are actually used (2%)
 - Consider 16 hierarchical levels, 64 compartments which gives 10²⁰ labels

EMBEDDING A POSET IN A LATTICE

- Smith's subset of 21 labels do form a lattice. In general, however, selecting a subset of labels from a given lattice
 - may not yield a lattice, but
 - is guaranteed to yield a partial ordering
- Given a partial ordering we can always add extra labels to make it a lattice

EMBEDDING A POSET IN A LATTICE





BLP BASIC ASSUMPTIONS

- SUB = {S1, S2, ..., Sm}, a fixed set of subjects
- OBJ = {O1, O2, ..., On}, a fixed set of objects
- $R \supseteq \{r, w\}$, a fixed set of rights
- D, an m \times n discretionary access matrix with D[i,j] \subseteq R
- M, an $m \times n$ current access matrix with $M[i,j] \subseteq \{r, w\}$

BLP MODEL (LIBERAL STAR-PROPERTY)

Lattice of confidentiality labels

$$\Lambda = {\lambda 1, \lambda 2, ..., \lambda p}$$

Static assignment of confidentiality labels

$$\lambda$$
: SUB \cup OBJ $\rightarrow \Lambda$

- M, an m × n current access matrix with
 - $r \in M[i,j] \Rightarrow r \in D[i,j] \land \lambda(Si) \ge \lambda(Oj)$ simple security
 - $w \in M[i,j] \Rightarrow w \in D[i,j] \land \lambda(Si) \le \lambda(Oj)$ star-property

BLP MODEL (STRICT STAR-PROPERTY)

Lattice of confidentiality labels

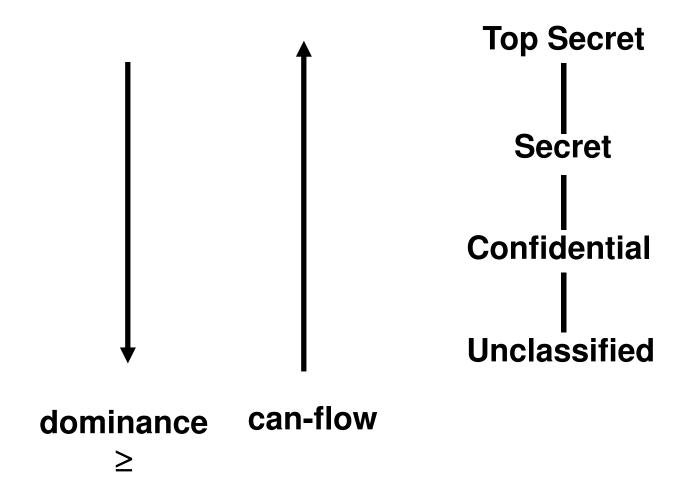
$$\Lambda = {\lambda 1, \lambda 2, ..., \lambda p}$$

Static assignment of confidentiality labels

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 - $w \in M[i,j] \Rightarrow w \in D[i,j] \land \lambda(Si) = \lambda(Oj)$ star-property

BLP MODEL



STAR-PROPERTY

- applies to subjects not to users
- users are trusted (must be trusted) not to disclose secret information outside of the computer system
- subjects are not trusted because they may have Trojan Horses embedded in the code they execute
- star-property prevents overt leakage of information and does not address the covert channel problem

BIBA MODEL

Lattice of integrity labels

$$\Omega = \{\omega 1, \omega 2, ..., \omega \mathbf{q}\}$$

Assignment of integrity labels

$$\omega$$
: SUB \cup OBJ $\rightarrow \Omega$

- M, an m × n current access matrix with
 - $r \in M[i,j] \Rightarrow r \in D[i,j] \land \omega(Si) \le \omega(Oj)$

simple integrity

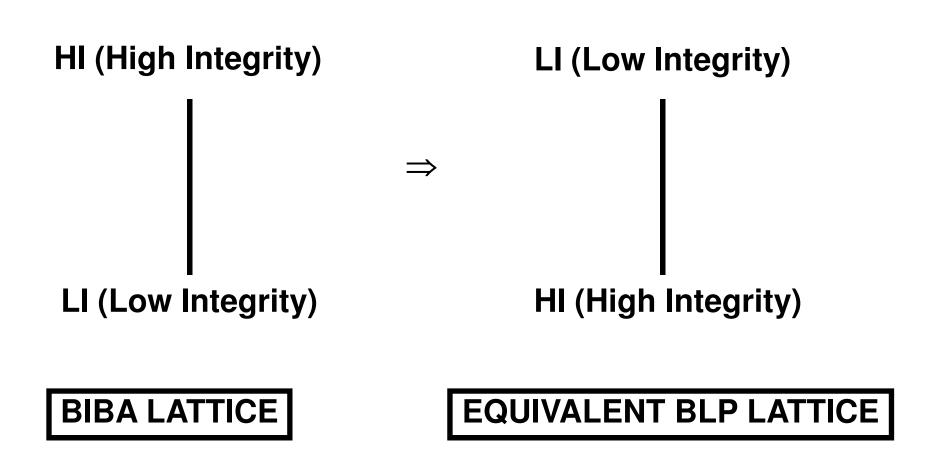
• $w \in M[i,j] \Rightarrow w \in D[i,j] \land \omega(Si) \ge \omega(Oj)$

integrity confinement

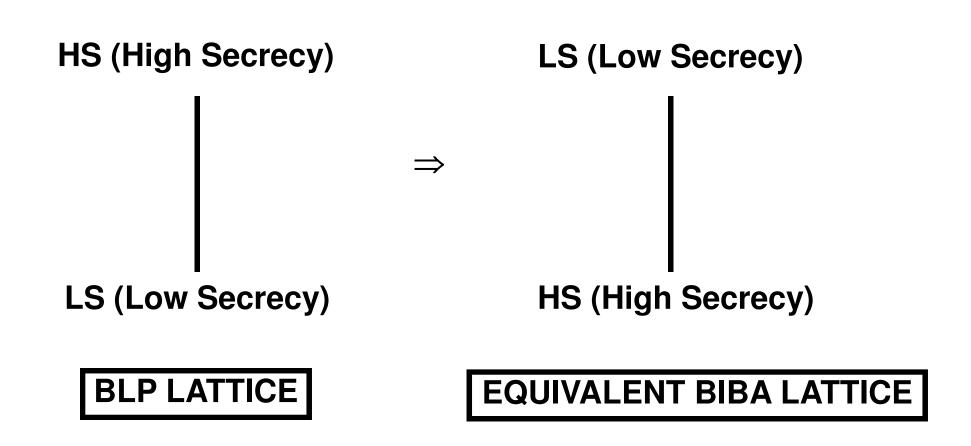
EQUIVALENCE OF BLP AND BIBA

- Information flow in the Biba model is from top to bottom
- Information flow in the BLP model is from bottom to top
- Since top and bottom are relative terms, the two models are fundamentally equivalent

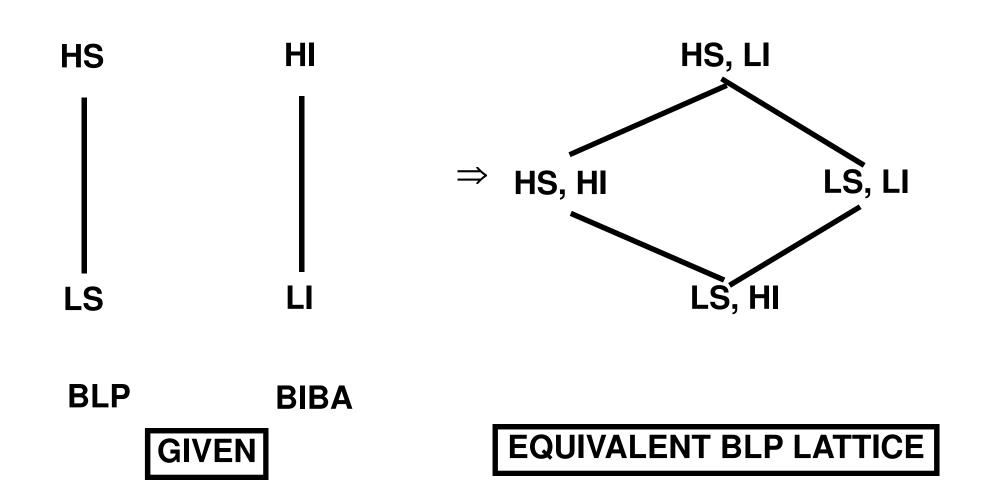
EQUIVALENCE OF BLP AND BIBA



EQUIVALENCE OF BLP AND BIBA

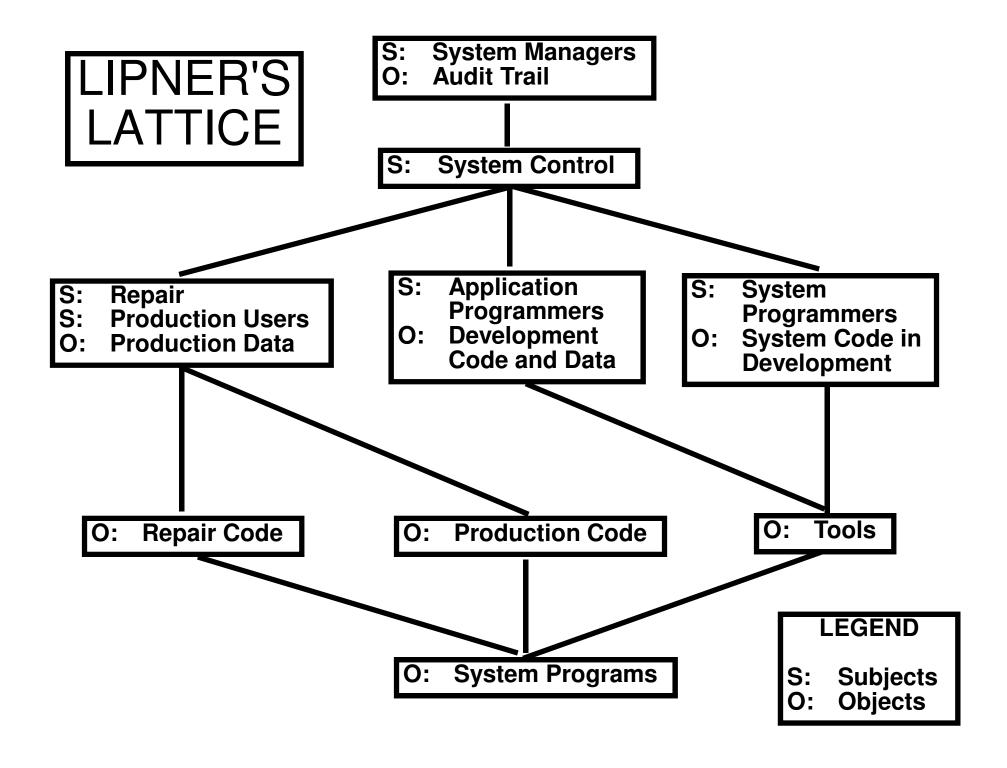


COMBINATION OF DISTINCT LATTICES



BLP AND BIBA

- BLP and Biba are fundamentally equivalent and interchangeable
- Lattice-based access control is a mechanism for enforcing one-way information flow, which can be applied to confidentiality or integrity goals
- We will use the BLP formulation with high confidentiality at the top of the lattice, and high integrity at the bottom



LIPNER'S LATTICE

- Lipner's lattice uses 9 labels from a possible space of 192 labels (3 integrity levels, 2 integrity compartments, 2 confidentiality levels, and 3 confidentiality compartments)
- The single lattice shown here can be constructed directly from first principles

LIPNER'S LATTICE

- The position of the audit trail at lowest integrity demonstrates the limitation of an information flow approach to integrity
- System control subjects are exempted from the star-property and allowed to
 - write down (with respect to confidentiality) or equivalently
 - write up (with respect to integrity)

DYNAMIC LABELS IN BLP

- Tranquility (most common):
 λ is static for subjects and objects
- BLP without tranquility may be secure or insecure depending upon the specific dynamics of labelling
- Noninterference can be used to prove the security of BLP with dynamic labels

DYNAMIC LABELS IN BLP

High water mark on subjects:

 λ is static for objects

 λ may increase but not decrease for subjects

Is secure and is useful

High water mark on objects:

 λ is static for subjects

 λ may increase but not decrease for subjects

Is insecure due to disappearing object signaling channel